CS4102 Algorithms

Spring 2020

Today's Keywords

- Greedy Algorithms
- Choice Function
- Cache Replacement
- Hardware & Algorithms

CLRS Reading: Chapter 16

Caching Problem

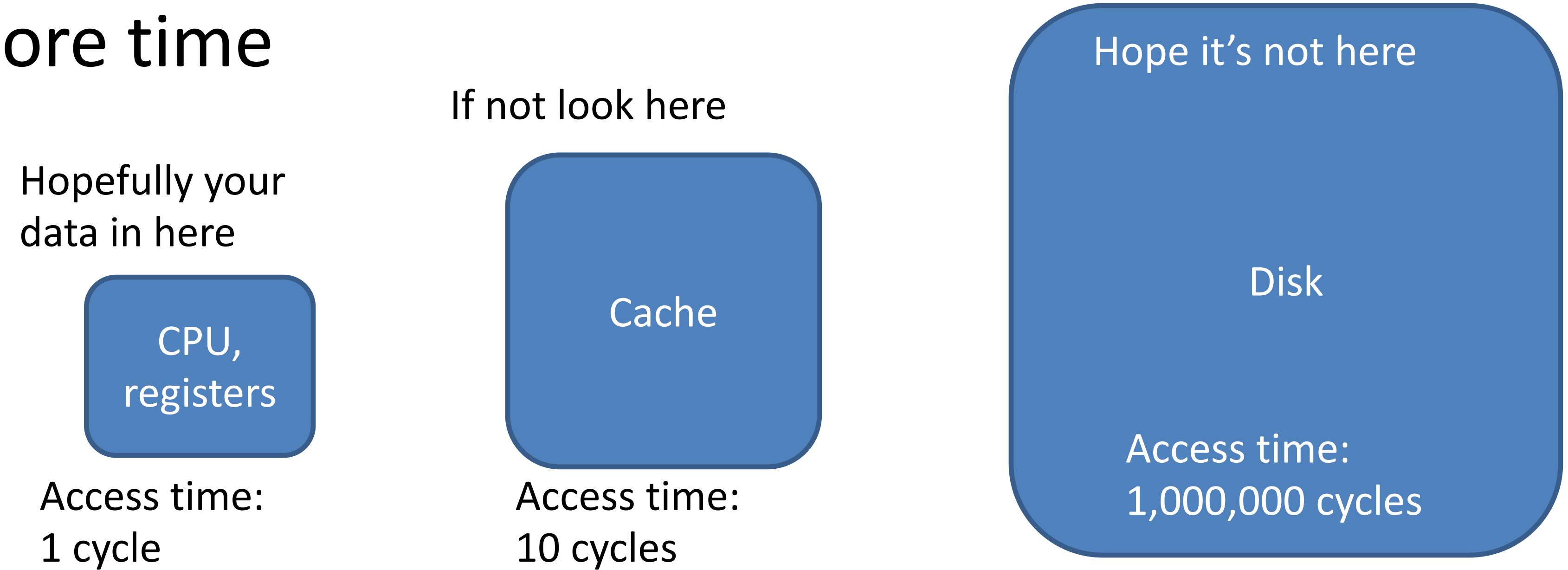
- Why is using too much memory a bad thing?

Von Neumann Bottleneck

- Named for John von Neumann
- Inventor of modern computer architecture
- Other notable influences include:
  - Mathematics
  - Physics
  - Economics
  - Computer Science

Von Neumann Bottleneck

- Reading from memory is VERY slow
- Big memory = slow memory
- Solution: hierarchical memory
- Takeaway for Algorithms: Memory is time, more memory is a lot more time



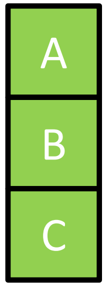
Caching Problem

- Cache misses are very expensive
- When we load something new into cache, we must eliminate something already there
- We want the best cache “schedule” to minimize the number of misses

Caching Problem Definition

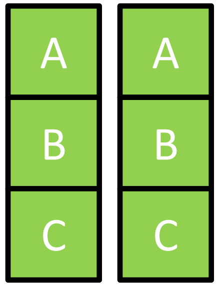
- Input:
  - k = size of the cache
  - $M = [m_1, m_2, \ldots m_n]$ = memory access pattern
- Output:
  - "schedule" for the cache (list of items in the cache at each time) which minimizes cache fetches

Example



A B C D A D E A D B A E C E A

Example



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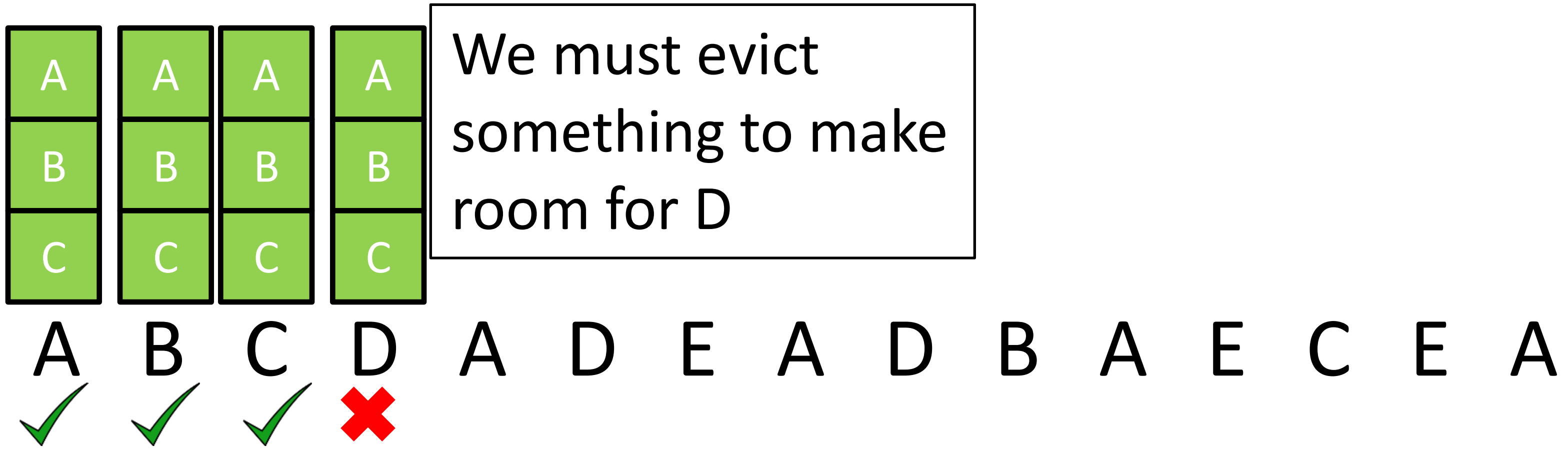
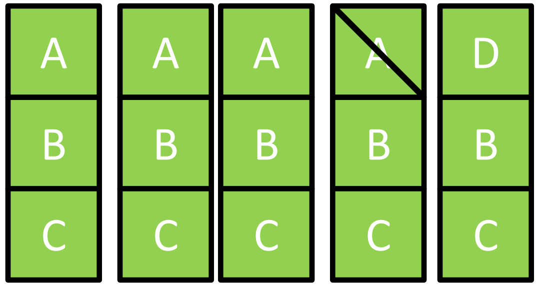

A B C
A B C
A B C
A B C
We must evict something to make room for D
A B C D A D E A D B A E C E A

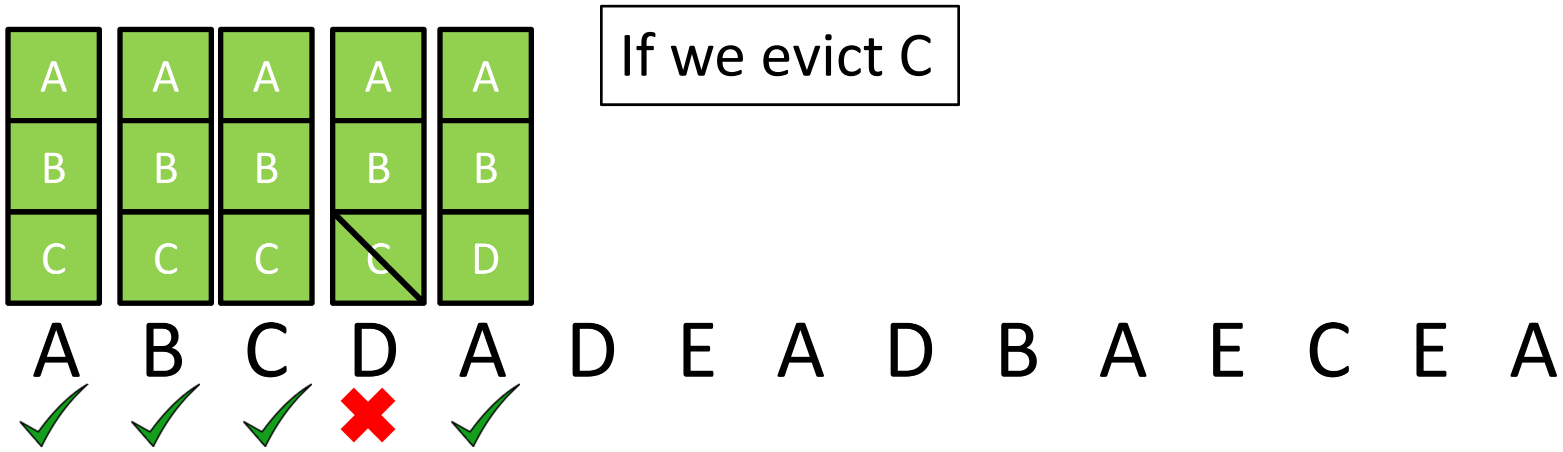
Example



If we evict A

A B C D A D E A D B A E C E A

Example



Our Problem vs Reality

- Assuming we know the entire access pattern
- Cache is Fully Associative
- Counting # of fetches (not necessarily misses)
- “Reduced” Schedule: Address only loaded on the cycle it’s required
  - Reduced == Unreduced (by number of fetches)

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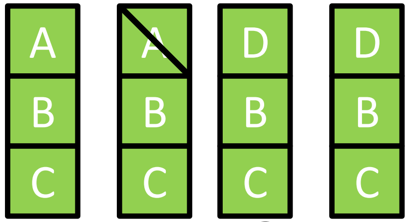
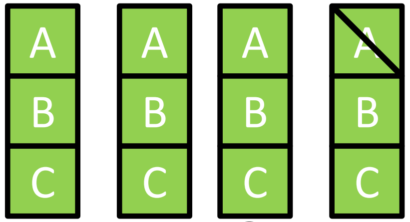


Unreduced

A B C D A D E A D B A E C E A



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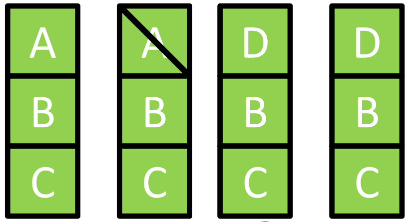

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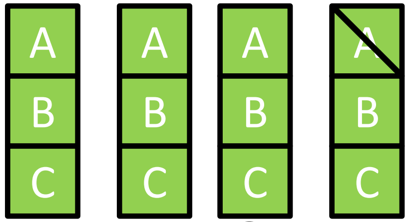

Reduced

Leaving A in longer does not save fetches

A B C D A D E A D B A E C E A

Greedy Algorithms

- Require Optimal Substructure
  - – Solution to larger problem contains the solution to a smaller one
  - – Only one subproblem to consider!
- Idea:
  1. Identify a greedy choice property
     - How to make a choice guaranteed to be included in some optimal solution
  2. Repeatedly apply the choice property until no subproblems remain

Greedy choice property

- Belady evict rule:
  - Evict the item accessed farthest in the future

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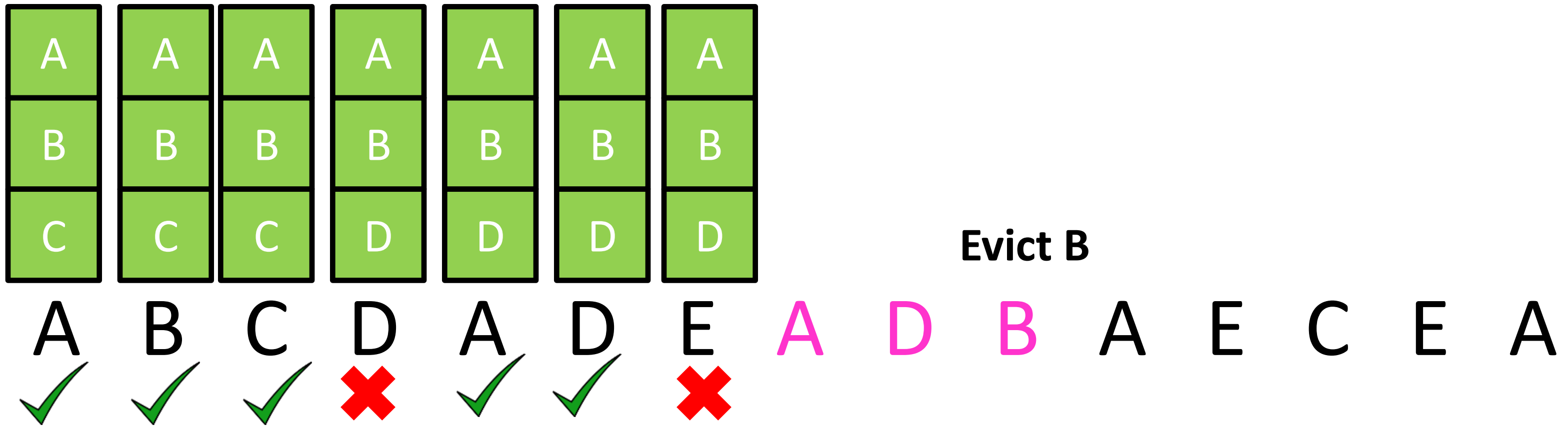

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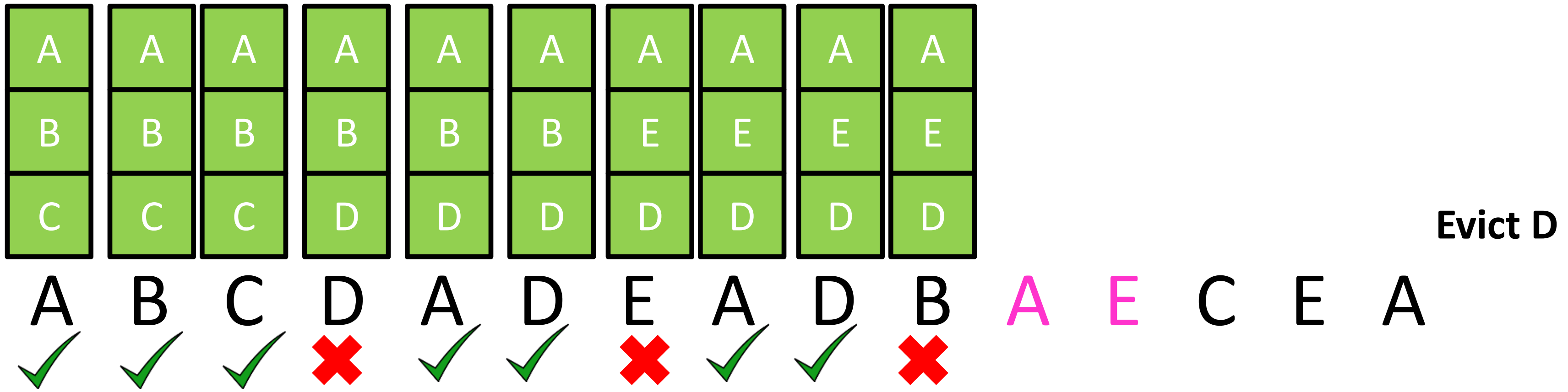

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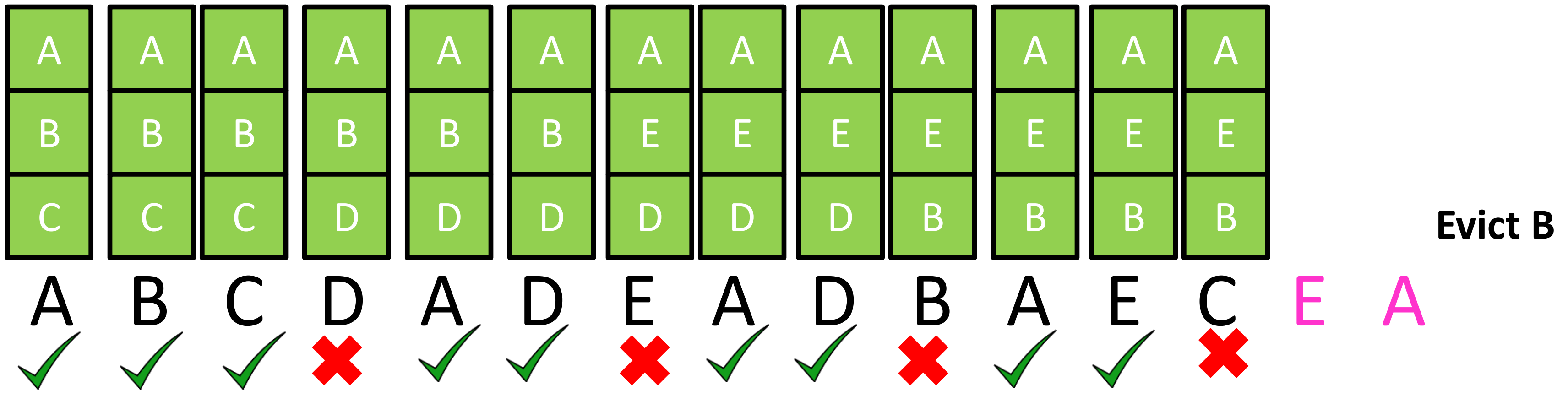

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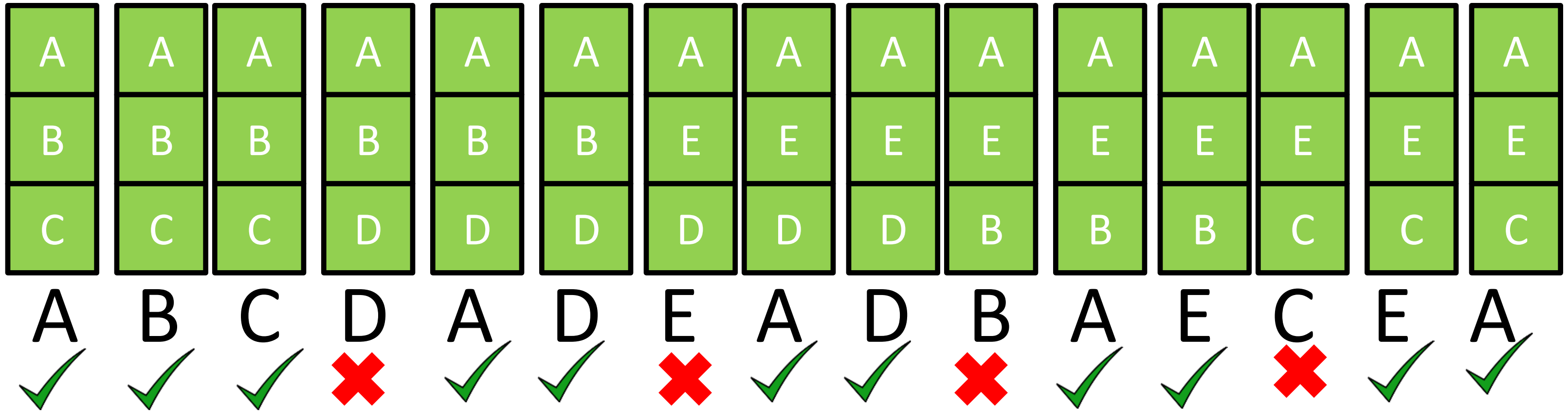

4 Cache Misses

Greedy Algorithms

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Caching Greedy Algorithm

Initialize $cache$= first k accesses

For each $m_i \in M$:

    if $m_i \in cache$:

        print $cache$

    else:

        $m =$ furthest-in-future from cache

        evict m, load m_i

        print $cache$

Exchange argument

- Shows correctness of a greedy algorithm
- Idea:
  - Show exchanging an item from an arbitrary optimal solution with your greedy choice makes the new solution no worse
  - How to show my sandwich is at least as good as yours:
    - Show: “I can remove any item from your sandwich, and it would be no worse by replacing it with the same item from my sandwich”

Belady Exchange Lemma

Let S_{ff} be the schedule chosen by our greedy algorithm

Let S_i be a schedule which agrees with S_{ff} for the first i memory accesses.

We will show: there is a schedule S_{i+1} which agrees with S_{ff} for the first $i+1$ memory accesses, and has no more misses than S_i

(i.e. $misses(S_{i+1}) \leq misses(S_i)$)

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Optimal

$\boxed{S^*}$

Agrees with S_{ff} on first 0 accesses

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Greedy

S_{ff}

Agrees with S_{ff} on all n accesses

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Agrees with S_{ff} on first 0 accesses

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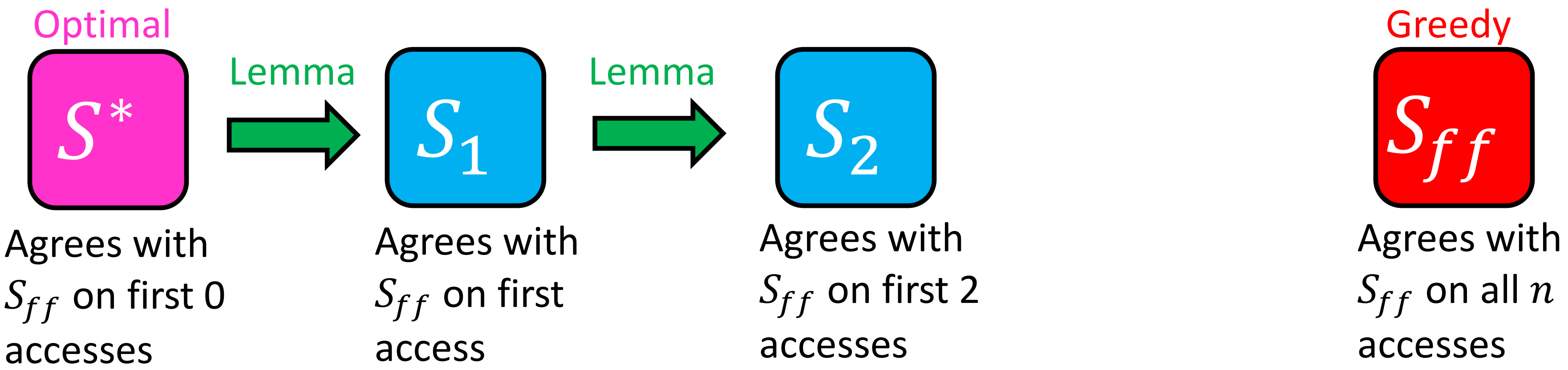

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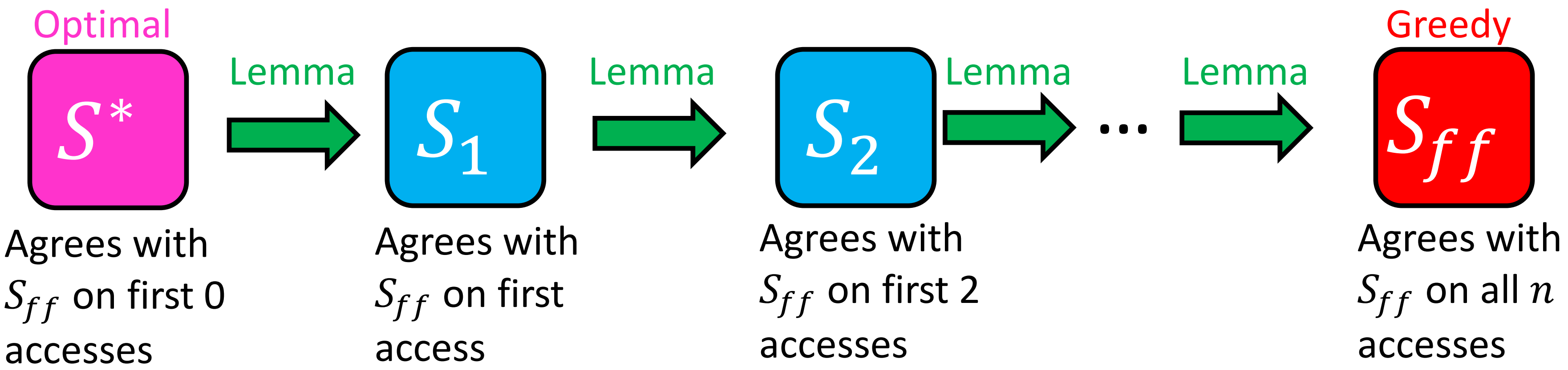

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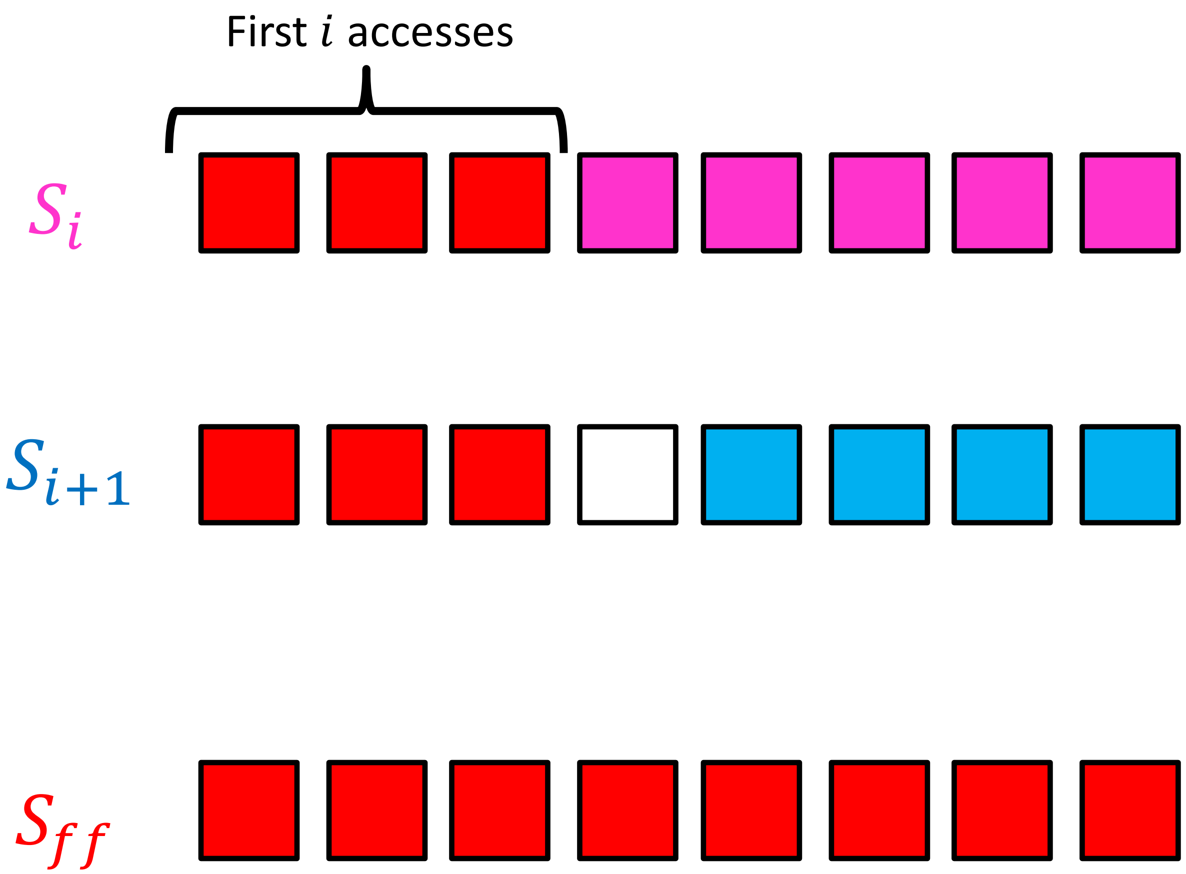

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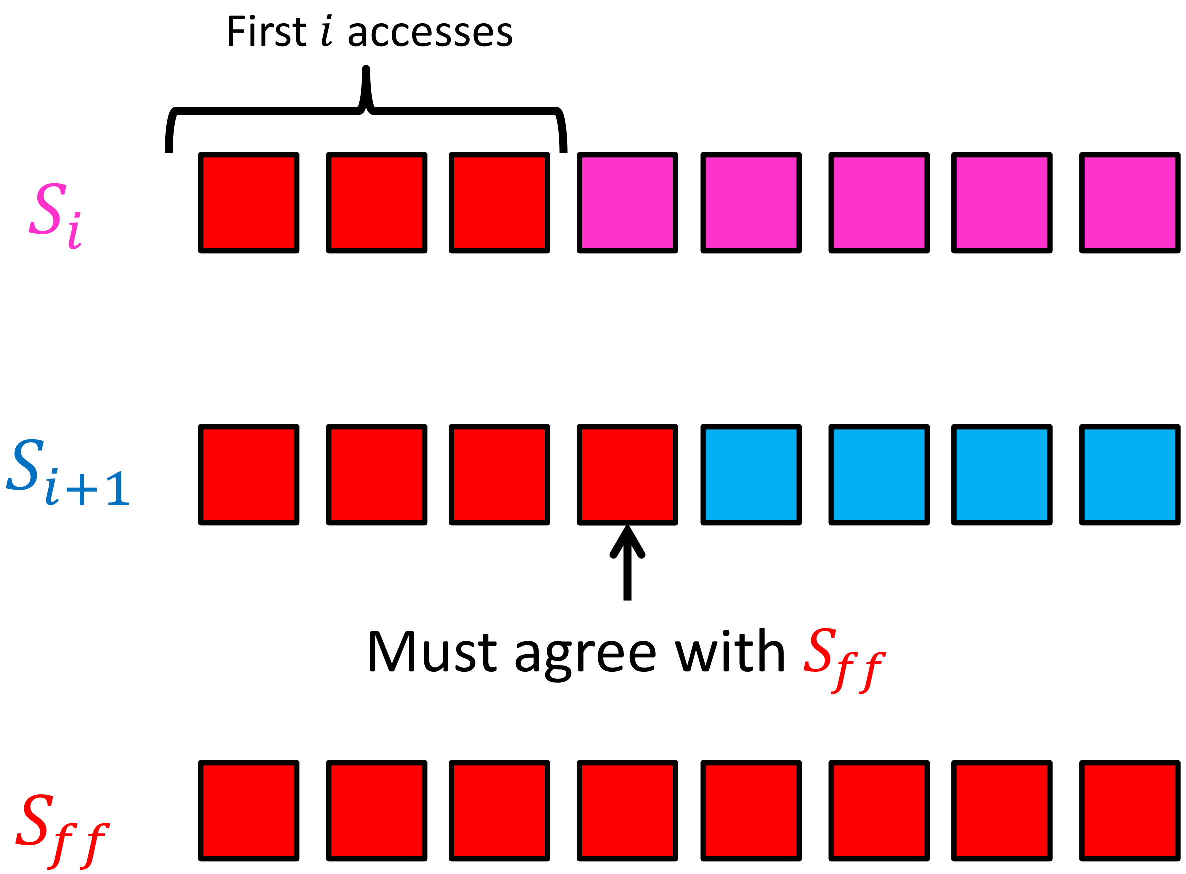
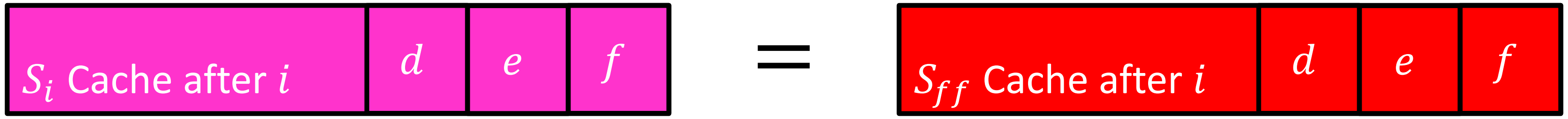


Need to fill in the rest of S_{i+1} to have no more misses than S_i

Proof of Lemma

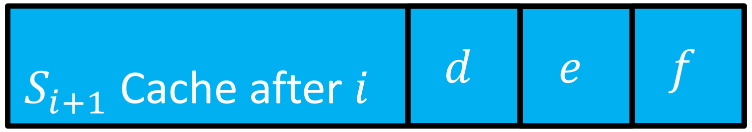
Goal: find S_{i+1} s.t. $misses(S_{i+1}) \leq misses(S_i)$

Since S_i agrees with S_{ff} for the first i accesses, the state of the cache at access $i+1$ will be the same



Consider access $m_{i+1} = d$

Case 1: if d is in the cache, then neither S_i nor S_{ff} evict from the cache, use the same cache for S_{i+1}



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S_i Cache after i	e	f

$=$

S_{ff} Cache after i	e	f

Consider access $m_{i+1} = d$

Case 2: if d isn't in the cache, and both S_i and S_{ff} evict f from the cache, evict f for d in S_{i+1}

S_{i+1} Cache after i	e	d

Proof of Lemma

Goal: find S_{i+1} s.t. $misses(S_{i+1}) \leq misses(S_i)$

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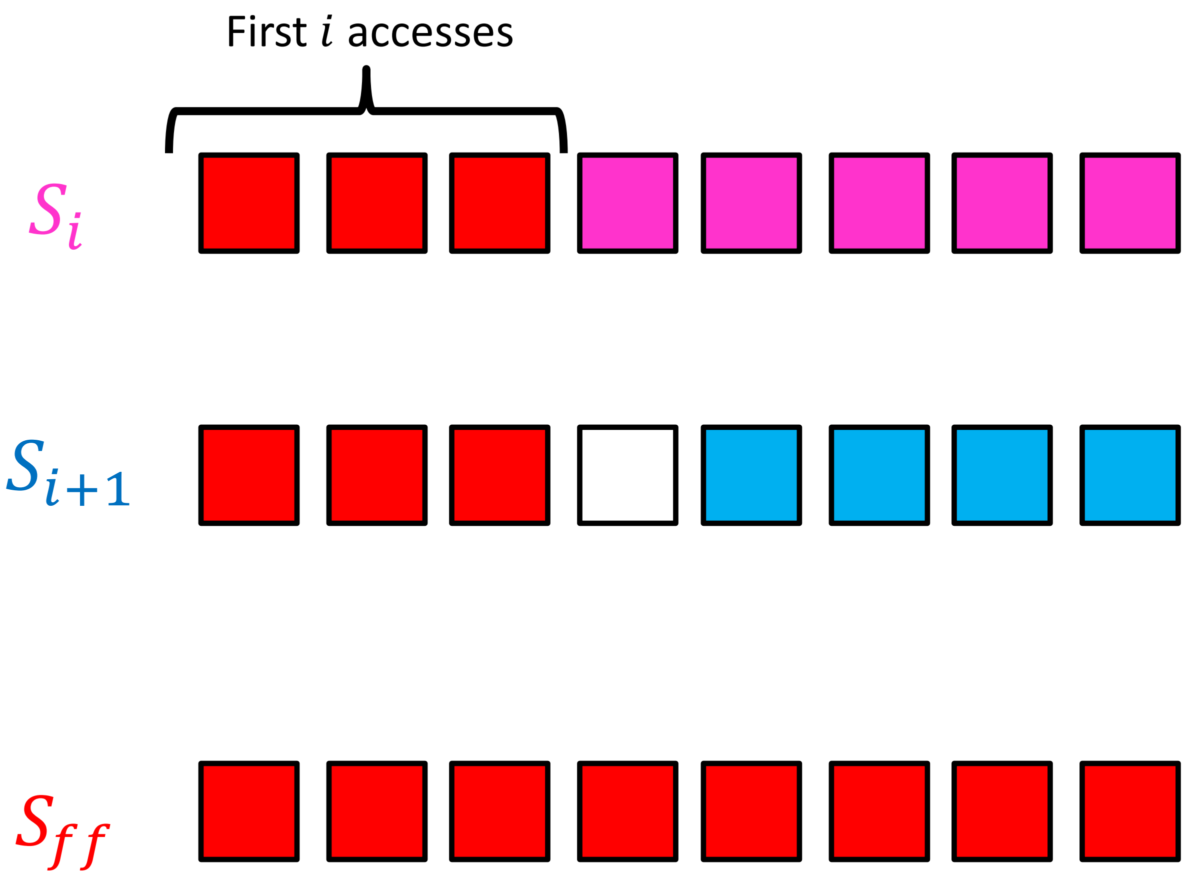
S_i Cache after i	e	f	$=$	S_{ff} Cache after i	e	f

Consider access $m_{i+1} = d$

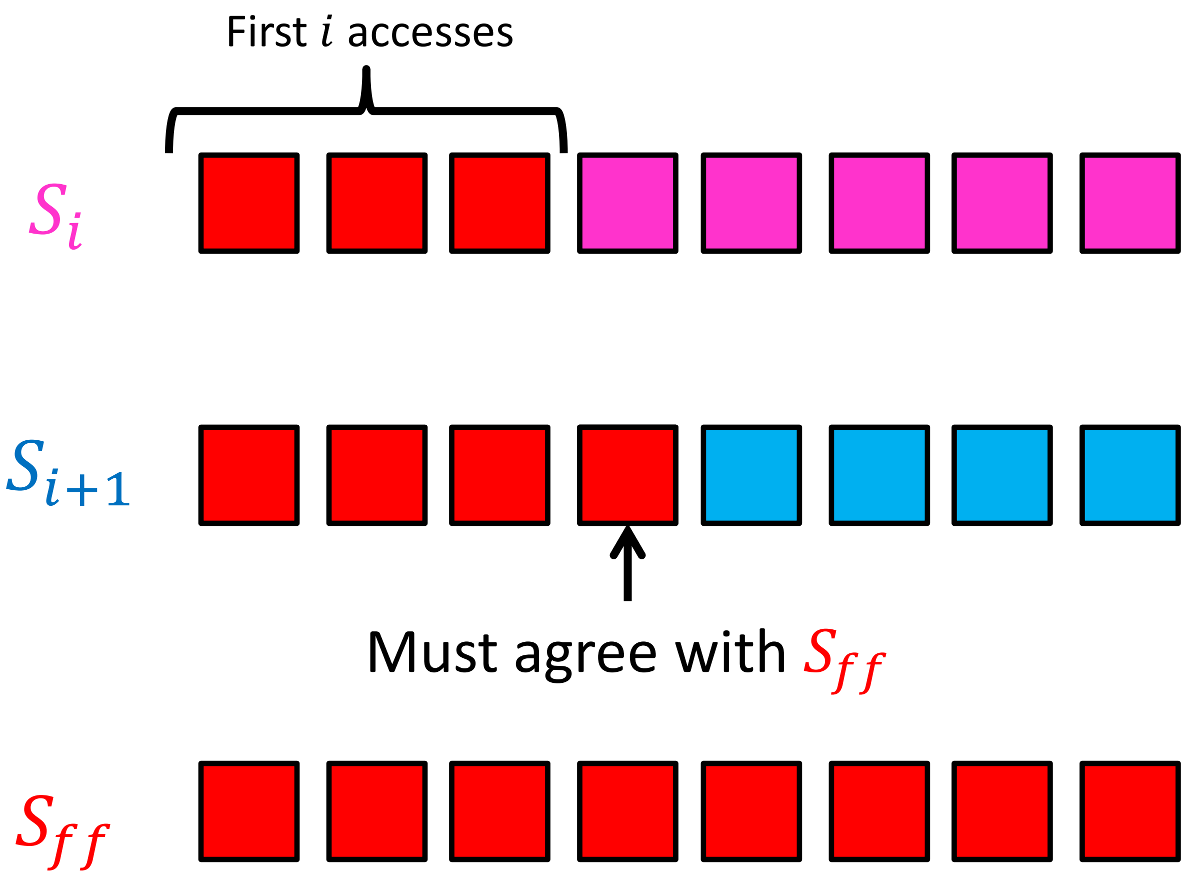
Case 3: if d isn't in the cache, S_i evicts e and S_{ff} evicts f from the cache

S_i Cache after $i+1$	d	f	$\neq$	S_{ff} Cache after $i+1$	e	d

Case 3



Case 3



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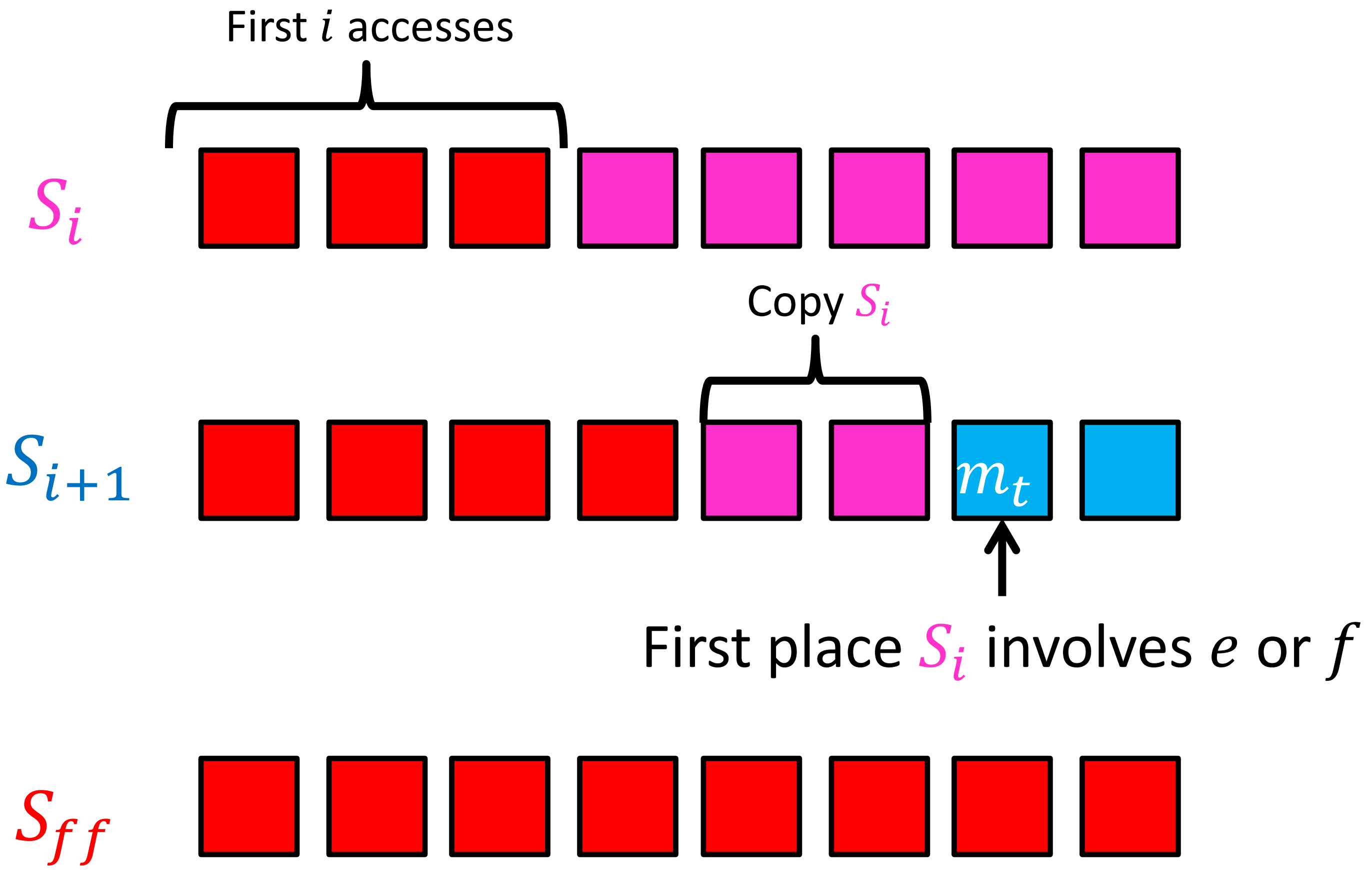
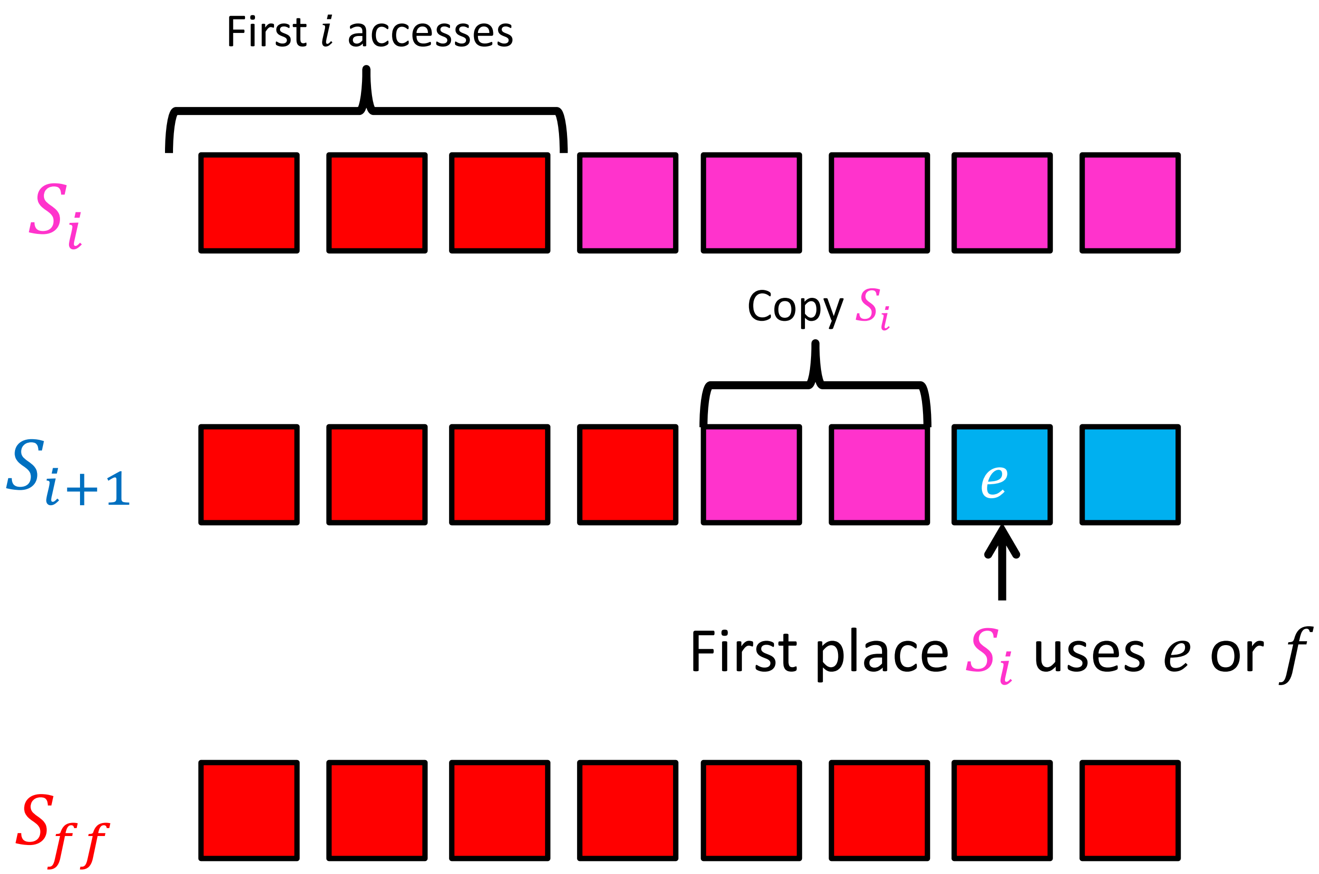


m_t = the first access after $i+1$ in which S_i deals with e or f

3 options: $\boldsymbol{m_t = e}$ or $\boldsymbol{m_t = f}$ or $\boldsymbol{m_t = x \neq e, f}$

Case 3, $m_t = e$



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Goal: find S_{i+1} s.t. $misses(S_{i+1}) \le misses(S_i)$

S_i Cache after $t-1$	x	d	f	$\neq$	S_{i+1} Cache after $t-1$	x	e	d

S_i must load e into the cache, assume it evicts x

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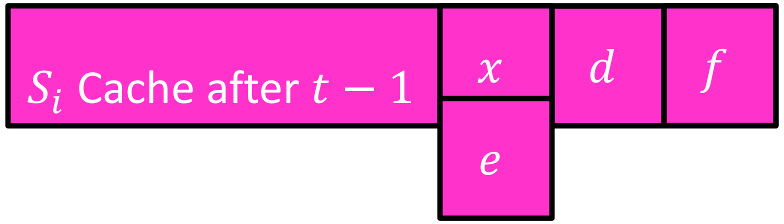
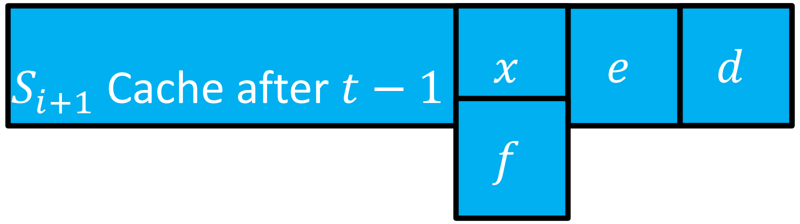


$\neq$



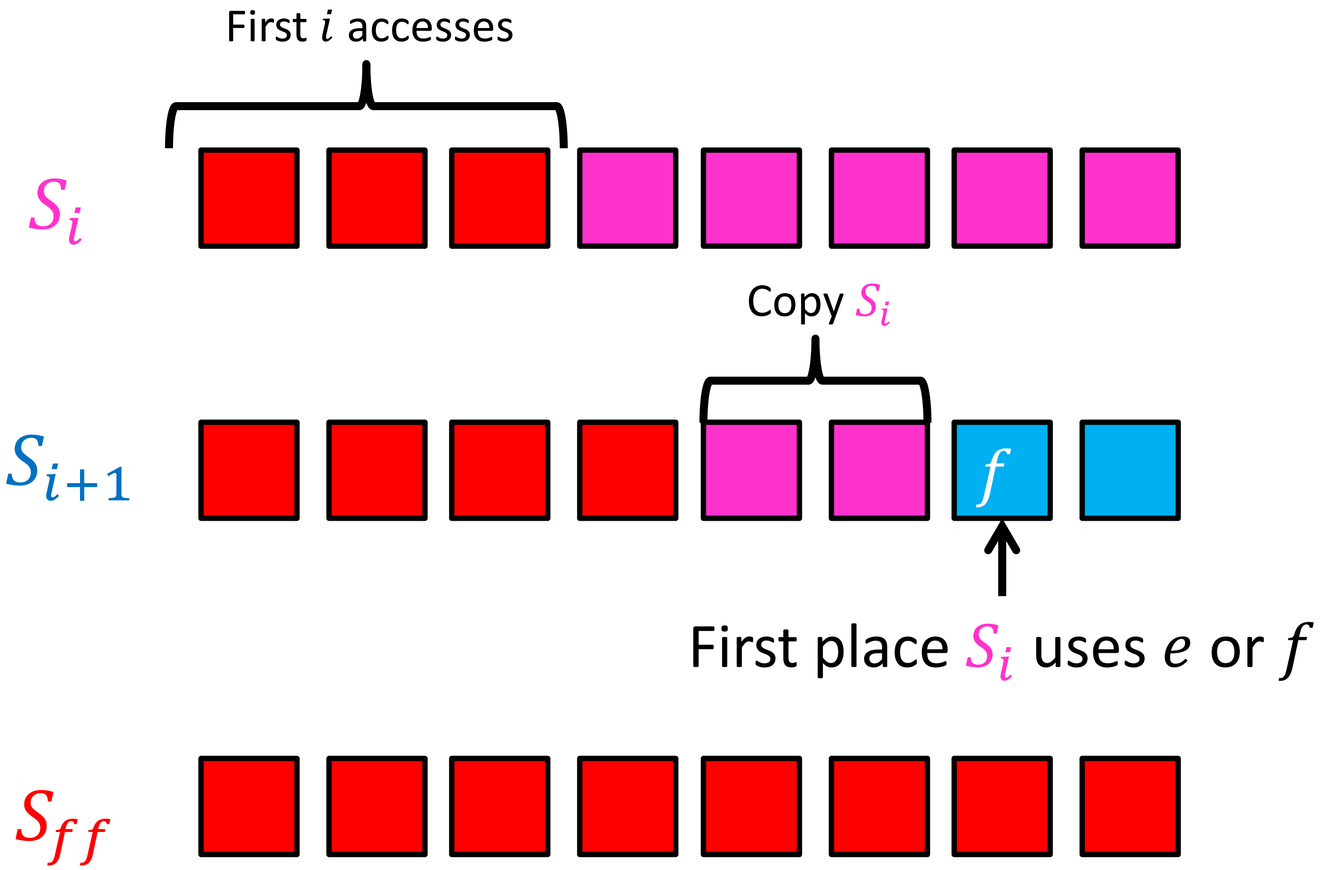
S_i must load e into the cache, assume it evicts x

S_{i+1} will load f into the cache, evicting x

The caches now match!

S_{i+1} behaved exactly the same as S_i between i and t, and has the same cache after t, therefore $misses(S_{i+1}) = misses(S_i)$

Case 3, $m_t = f$



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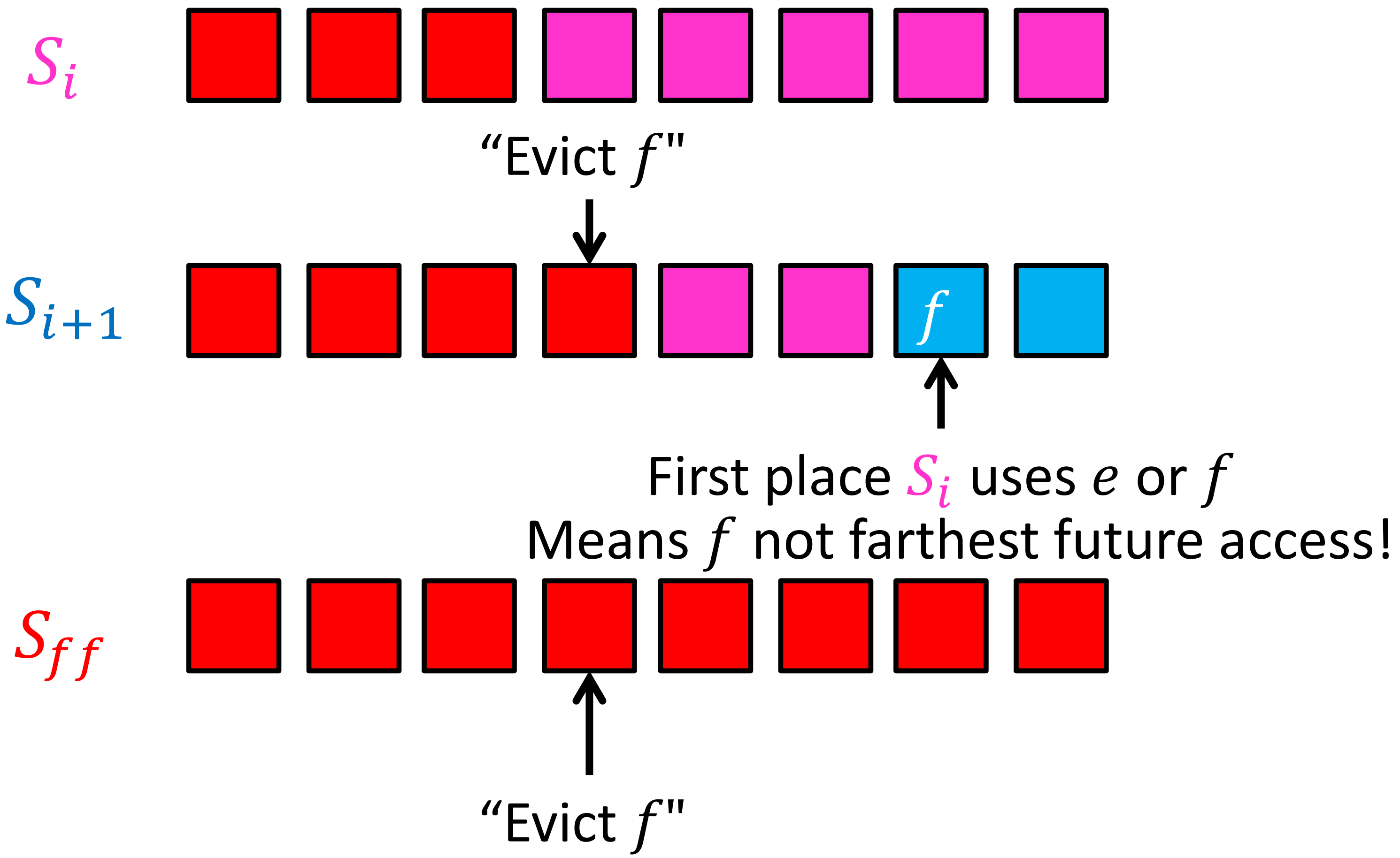
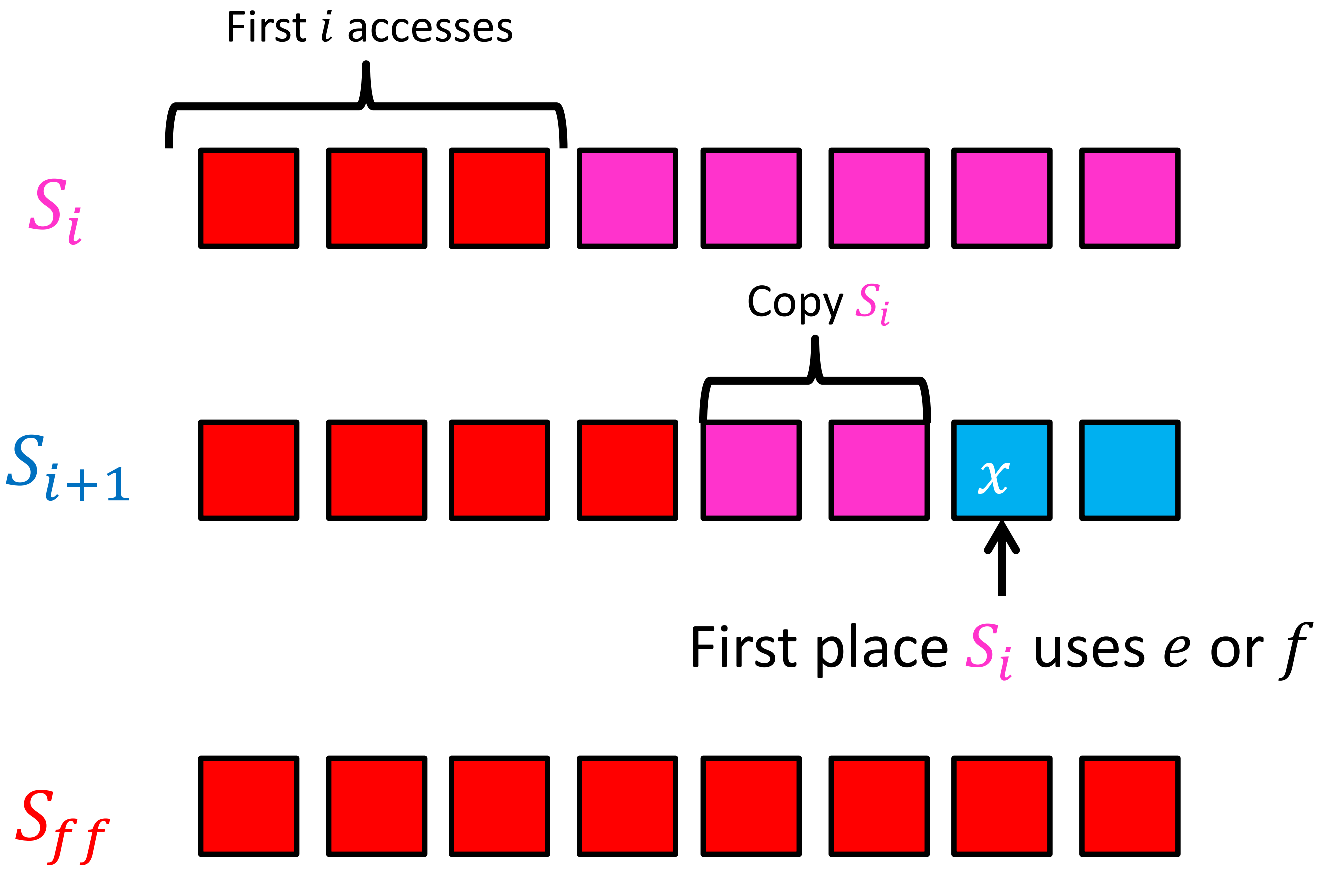
Cannot Happen!



Case 3, $m_t = x \neq e, f$



m_t = the first access after $i + 1$ in which S_i deals with e or f

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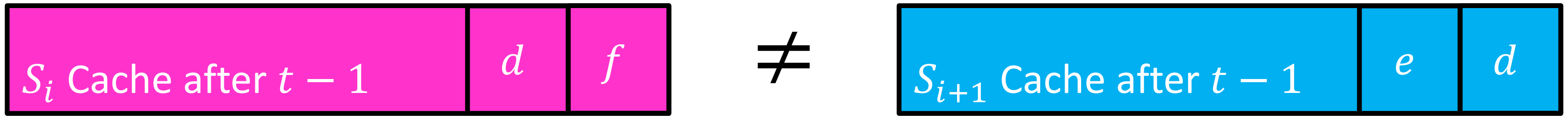

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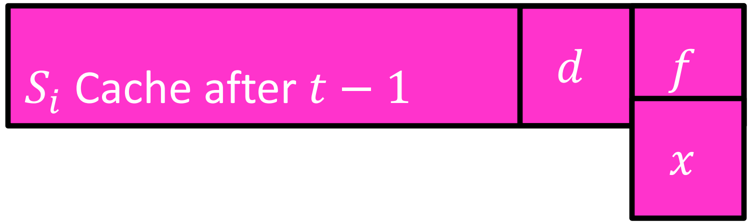

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Use Lemma to show Optimality

$$S^* \xrightarrow{\text{Lemma}} S_1 \xrightarrow{\text{Lemma}} S_2 \xrightarrow{\text{Lemma}} \cdots \xrightarrow{\text{Lemma}} S_{ff}$$

- S^*: Agrees with S_{ff} on first 0 accesses
- S_1: Agrees with S_{ff} on first access
- S_2: Agrees with S_{ff} on first 2 accesses
- S_{ff}: Agrees with S_{ff} on all n accesses